

# Posets

## A First Look

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# A Quick Introduction...

## For those of you I haven't met before...

Here's a few things about me:

- ▶ Hi! I'm Pranav!
- ▶ I'll be a freshman in college this fall.
- ▶ My favorite math subject is combo!
- ▶ Check me out at <https://pranavkonda.com>!

But that's enough about me, what you're really here for is the math.

## Some background about posets and notation

- ▶ Today we'll be learning about *partially ordered sets*.
- ▶ We also call them posets for short.
- ▶ They have cool connections to other areas of combo, like PIE.
- ▶ Since there's going to be some slightly advanced theory, let me know if I need to slow down or explain something in more detail.

The set  $\{1, 2, \dots, n\}$  is going to be denoted by  $[n]$ .

# Basic Definitions

# Binary Relations

As you probably already know, the *Cartesian Product* of two sets  $A$  and  $B$ ,  $A \times B$ , is the set

$$\{(a, b) : a \in A, b \in B\}.$$

## Definition (Binary Relation)

A binary relation  $R$  over two sets  $A$  and  $B$  is a subset of  $A \times B$ . We think about  $(a, b)$  as  $aRb$ , where we're *relating*  $a$  to  $b$ .

# Posets

Now we're ready to define a poset! A poset is a set  $P$ , along with a binary relation  $\leq$  that satisfies the following properties:

- ▶ **Reflexivity:** For any  $s \in P$ ,  $s \leq s$ .
- ▶ **Antisymmetry:** If  $t \leq s$  and  $s \leq t$ , then  $s = t$ .
- ▶ **Transitivity:** If  $s \leq u$  and  $u \leq t$ , then  $s \leq t$ .

Just like regular less-than,  $s < t$  indicates that  $s \neq t$ . Same for greater than signs.

# Posets

The “crux” of what makes the poset a *partially ordered* set is that for any two elements in  $P$ , it’s not necessary for  $s \leq t$  or  $t \leq s$  to hold! We have formal words for this (let  $s, t \in P$ ):

- ▶ If  $s \leq t$  or  $t \leq s$ , then we say  $s$  and  $t$  are **comparable**.
- ▶ Otherwise, we say  $s$  and  $t$  are **incomparable**.

Sometimes, our set and partial order will only have comparable elements, in which case we call it a total order (like  $(\mathbb{R}, \leq)$  for example).



# The $n$ chain

- ▶ Suppose  $n$  is a positive integer. Then the set  $[n]$  with the regular less-than relation forms a poset,  $\mathbf{n}$ , which we call the  $n$ -**chain**. It's pretty easy to verify that this satisfies reflexivity, antisymmetry, and transitivity.
- ▶ We'll see why it's called a *chain* shortly, but we need to define a few things first.

# Power sets

Your turn to prove something!

## Problem

*Let  $S$  be a set. Prove that  $(\mathcal{P}(S), \subseteq)$  is a poset.*

# Isomorphism

As expected with any mathematical object, we have some notion of poset *isomorphism*, or an *order-preserving bijection*. We say two posets  $P$  and  $Q$  are isomorphic if there's some map  $\phi : P \rightarrow Q$  such that

$$s \leq t \iff \phi(s) \leq \phi(t).$$

Pretty generic definition of isomorphism.

# Subsets

Subsets are a bit wonky, since there are two ways we can define one. Since we will almost always use only one kind, we're going to go through that.

## Definition (Induced subposet)

Let  $P$  be a poset. Then  $Q$  is an induced subposet of  $P$  if for  $s, t \in Q$ ,  $s \leq_Q t$  if and only if  $s \leq_P t$ . Obviously,  $Q \subseteq P$ .

## Covering relations

This is an important definition. For two elements  $s$  and  $t$  in  $P$ , we say that  $t$  *covers*  $s$  if  $s < t$  and there does not exist a  $u \in P$  such that  $s < u < t$ . We denote these *covering relations* by  $s \triangleleft t$ .

Covering relations are important because they let us draw posets! Formally, we call these drawings *Hasse diagrams*, and they're basically directed acyclic graphs but we understand every edge points "up".

## The $n$ -chain returns

Let's consider the  $n$ -chain again, by drawing its Hasse diagram for  $n = 4$ . Obviously,  $n < n + 1$ , so we get something like this:



A lot of posets look really cool when we draw them.

## Power sets again

Your turn again! Try drawing the Hasse diagram of  $(\{1, 2, 3\}, \subseteq)$ . We also have a formal term for the poset  $([n], \subseteq)$ : the *Boolean poset*,  $B_n$ . Be sure to get familiar with this poset, as it's going to pop up again and again.

# Chains and Generating Functions



## Chains, in general

- ▶ Chains are actually a more general component of posets.
- ▶ A chain is a poset that is totally ordered: every element is comparable. So  $\mathbb{R}$  and  $\mathbb{Z}$  are chains.
- ▶ A subset  $C \subset P$  is a chain if every element in it is comparable, this is the more useful definition.
- ▶ We also have a notion of *length*: the length of a chain  $C \subset P$  is  $\ell(C) = |C| - 1$ .

## Rank and graded posets

The *rank* of a poset  $P$  is the length of its maximal chain:

$$\text{rank } P = \max_{C \subset P} \ell(C).$$

When every maximal chain (meaning there isn't a larger chain that contains it) of  $P$  has the *same* length, we call  $P$  a *graded* poset. If that length is  $n$ , then  $P$  is graded of rank  $n$ .

## Rank functions

If  $P$  is a graded poset of rank  $n$ , then there exists a *rank function*  $\rho : P \rightarrow \{1, 2, \dots, n\}$ , which satisfies these special properties:

- ▶ If  $s$  is a minimal element of  $P$ , then  $\rho(s) = 0$ .
- ▶ If  $s \lessdot t$ , then

$$\rho(s) + 1 = \rho(t).$$

If  $\rho(s) = i$ , then we say  $s$  has rank  $i$ .

## Rank generating functions

Suppose that  $P$  is a graded poset of rank  $n$ . Then there exists a *rank generating function*

$$F(P, x) = \sum_{k=0}^n p_k x^k,$$

where  $p_k$  is the number of elements in  $P$  with rank  $k$ .

## Rank generating functions

Can you try finding the rank function and rank generating function for the following posets?

- ▶ The  $n$ -chain  $\mathbf{n}$ .
- ▶ The boolean poset  $B_n$ .
- ▶ This is the other major poset I wanted to talk about: the *Divisor poset*  $D_n$ . For this poset, we take a positive integer  $n$ , and let the ground set be the positive integer divisors of  $n$ . Then  $s \leq t$  if  $s|t$ .

# Antichains

- ▶ An *antichain* is a subset  $A \subset P$  such that *no* two elements of  $A$  are comparable.
- ▶ Antichains are also called **Sperner families**, as they have a very deep connection to **Sperner's theorem** in extremal combinatorics.
- ▶ We'll look at some results related to extremal combinatorics (Dilworth's and Mirsky's theorems) some other time.

# Order Ideals

## Definition (Order Ideals)

An **order ideal** (also called a down set) is a subset  $I \subset P$  such that if  $t \in I$  and  $s \leq t$ ,  $s \in I$ .

- ▶ There's a *dual* of an order-ideal, called an *up-set*, where we replace the  $s \leq t$  condition with  $s \geq t$ .
- ▶ For experts, this should remind you of the definition of a regular ideal.

## A nice problem

### Problem

*Show that the number of order ideals in  $P$  is equal to the number of antichains in  $P$ .*

Hint: Argue for a bijection!



## References and places to learn more

- ▶ *R. Stanley, Enumerative Combinatorics Volume 1.*
- ▶ One of the best combo textbooks out there, although it's very dense and reference-like.
- ▶ More of these lectures!